## Lecture 7: <br> Mesh representations and Mesh Processing

Computer Graphics: Rendering, Geometry, and Image Manipulation Stanford CS248A, Winter 2024

## A small triangle mesh



8 vertices, 12 triangles

## A large triangle mesh

## David

Digital Michelangelo Project
28,184,526 vertices
56,230,343 triangles


## Even larger meshes

## Google Farth

Meshes reconstructed from satellite and aertal photography
Trillions of triangles

## Recall: image upsampling



Convert representation of signal given by samples taken at black dots into a representation given at new set of denser samples (red dots)

## Recall: image upsampling



## Recall: image upsampling



## Recall: image downsampling



Convert representation of signal given by samples taken at black dots into a representation given at new set of sparser samples (red dots)

## Recall: image resampling



Convert representation of signal given̂ by samples taken at black dots into a representation given at new set of samples (red dots)

## Examples of geometry processing

## Mesh upsampling — subdivision



Increase resolution via interpolation

## Mesh downsampling — simplification



Decrease resolution; try to preserve shape/appearance

## Mesh resampling - regularization



Modify sample distribution to improve quality

## More geometry processing tasks



## Today

- How to represent meshes (data structures)
- How to perform a number of basic mesh processing operations
- Subdivision (upsampling)
- Mesh simplification (downsampling)
- Mesh resampling


## Mesh representations

## Basic mesh representation: list of triangles



## Another representation: <br> Lists of vertexes / indexed triangle



## Comparison

- List of triangles
- GOOD: simple
- BAD: contains redundant per-vertex information
- List of vertexes + list of indexed triangles
- GOOD: sharing vertex position information reduces memory usage
- GOOD: ensures integrity of the mesh (changing a vertex's position in 3D space causes that vertex in all the polygons to move)


## Mesh topology vs surface geometry

Same vertex positions, different mesh topology


Same topology, different vertex positions


## Smooth surfaces

- Intuitively, a surface is the boundary or "shell" of an object
- (Think about the candy shell, not the chocolate.)
- Surfaces are manifold:
- If you zoom in far enough (at any point) looks like a plane*
- E.g., the Earth from space vs. from the ground



## Why is the manifold property valuable?

- Makes life simple: all surfaces look the same (at least locally)
- Gives us coordinates! (at least locally)



## Isn't every shape manifold?

- No, for instance:


Center point never looks like the plane, no matter how close we get.

## More examples of smooth surfaces

- Which of these shapes are manifold?



## A manifold polygon mesh has fans, not fins

- For polygonal surfaces just two easy conditions to check:

1. Every edge is contained in only two polygons (no"fins")
2. The polygons containing each vertex make a single "fan"


## What about boundary?

- The boundary is where the surface "ends."
- E.g., waist and ankles on a pair of pants.
- Locally, looks like a half disk

■ Globally, each boundary forms a loop

■ Polygon mesh:


- one polygon per boundary edge
- boundary vertex looks like"pacman"


## Topological validity: manifold

A 2D manifold is a surface that when cut with a small sphere always yields a disk (or a half disk on the boundary)

Manifold


With border


Not manifold


With border


## Manifolds have useful properties

- A 2D manifold is a surface that when cut with a small sphere always yields a disk
- If a mesh is manifold, we can rely on these useful properties: *
- An edge connects exactly two faces
- An edge connects exactly two vertices
- A face consists of a ring of edges and vertices
- A vertex consists of a ring of edges and faces
- Euler's polyhedron formula holds: \#f - \#e + \#v = 2 (for a surface topologically equivalent to a sphere) (Check for a cube: $6-12+8=2$ )


## Topological validity: orientation consistency

Both facing front


Inconsistent orientations


Non-orientable
(e.g., Moebius strip)



## Simple example: triangle-neighbor data structure

```
// definition of a triangle
struct Tri {
    Vert* v[3];
    Tri* t[3];
}
// definition of a triangle vertex
struct Vert {
    Vec3 pos;
    Tri* t;
}
```



## Triangle-neighbor - mesh traversal

Find next triangle counter-clockwise around vertex v from triangle $t$

```
Tri* ccw_tri(Vert *v, Tri *t)
{
    if (v == t->v[0])
        return t[0];
    if (v == t->v[1])
        return t[1];
    if (v == t->v[2])
        return t[2];
}
```



## Half-edge data structure

```
struct Halfedge {
    Halfedge *twin,
    Halfedge *next;
    Vertex *vertex;
    Edge *edge;
    Face *face;
}
struct Vertex {
    Vec3 pos;
    Halfedge *halfedge;
}
struct Edge {
    Halfedge *halfedge;
}
struct Face {
    Halfedge *halfedge;
}
```

Key idea: two half-edges act as "glue" between mesh elements


## Each vertex, edge and face points to one of its half edges

## Half-edge structure facilitates mesh traversal

- Use twin and next pointers to move around mesh
- Process vertex, edge, and/or face pointers

```
Example 1: process all vertices of a face
Halfedge* h = f->halfedge;
do {
    do_work(h->vertex);
    h = h->next;
}
while( h != f->halfedge );
```



## Half-edge structure facilitates mesh traversal

## Example 2: process all edges around a vertex

```
Halfedge* h = v->halfedge;
do {
    do_work(h->edge);
    h = h->twin->next;
}
while( h != v->halfedge );
```



## Local mesh operations

## Half-Edge - local mesh editing

■ Consider basic operations for linked list: insert, delete

- Basic ops for half-edge mesh: flip, split, collapse edges


Allocate / delete elements; reassign pointers
(Care is needed to preserve mesh manifold property)

## Half-edge - edge flip

Triangles (a,b,c), (b,d,c) become (a,d,c), (a,b,d):


- In implementaton: you'll perform a long list of half-edge pointer reassignments
- However, no mesh elements created/destroyed


## Thought experiment: defining edge flip on N -gons?

- I find it very use to think about this case...
- What is a "reasonable" thing to do.
- Does your approach reduce to triangle edge flips in the $\mathrm{N}=3$ case?



## Half-edge - edge split

Insert midpoint m of edge ( $\mathbf{c}, \mathrm{b}$ ), connect to get four triangles:


- Must add elements to mesh (new vertex, faces, edges)
- Again, many half-edge pointer reassignments


## Half-edge - edge collapse

Replace edge ( $\mathrm{c}, \mathrm{d}$ ) with a single vertex m :


- Must delete elements from the mesh
- Again, many half-edge pointer reassignments


## Global mesh operations: geometry processing

- Mesh subdivision (form of subsampling)
- Mesh simplification (form of downsampling)
- Mesh regularization (form of resampling)



## Subdivision — upsampling a mesh

## Upsampling via subdivision



- Replace vertex positions with weighted average of neighbors
- Main considerations:
- interpolating vs. approximating
- limit surface continuity ( $\left.C_{1}, C_{2}, \ldots.\right)$
- behavior at irregular vertices

- Many options:
- Quad: Catmull-Clark
- Triangle: Loop, butterfly, sqrt(3)



## Loop subdivision

Common subdivision rule for triangle meshes "C2" smoothness away from irregular vertices
Approximating, not interpolating

uemxyn」 uow!S

## Loop subdivision algorithm

- Split each triangle into four

- Compute new vertex positions using weighted sum of prior vertex positions:

$\mathrm{n}=$ vertex degree
$u=3 / 16$ if $n=3,3 /(8 n)$ otherwise

New vertices
(weighted sum of vertices on split edge, and vertices "across from" edge)

Old vertices
(weighted sum of edge adjacent vertices)

## Loop subdivision algorithm

## Example, for degree 6 vertices ("regular" vertices)



## Loop subdivision results

Common subdivision rule for triangle meshes
"C2" smoothness away from irregular vertices
Approximating, not interpolating


## Semi-regular meshes

Most of the mesh has vertices with degree 6

But if the mesh is topologically equivalent to a sphere, then not all the vertices can have degree 6

Must have a few extraordinary points (degree not equal to 6)


## Proof: always an extraordinary vertex

Our triangle mesh (topologically equivalent to sphere) has V vertices, E edges, and T triangles
$E=3 / 2 \mathrm{~T}$

- There are 3 edges per triangle, and each edge is part of 2 triangles
- Therefore E = 3/2T
$T=2 V-4$
- Euler Convex Polyhedron Formula: T - E + V = 2
- $\Rightarrow \mathrm{V}=3 / 2 \mathrm{~T}-\mathrm{T}+2 \Rightarrow \mathrm{~T}=2 \mathrm{~V}-4$

If all vertices had 6 triangles, $\mathrm{T}=\mathbf{2 V}$

- There are 6 edges per vertex, and every edge connects 2 vertices
- Therefore, $\mathrm{E}=6 / 2 \mathrm{~V} \Rightarrow 3 / 2 \mathrm{~T}=6 / 2 \mathrm{~V} \Rightarrow \mathrm{~T}=2 \mathrm{~V}$

T cannot equal both 2V-4 and 2V, a contradiction

- Therefore, the mesh cannot have 6 triangles for every vertex


## Loop subdivision via edge operations

First, split edges of original mesh in any order:


Next, flip new edges that touch a new and old vertex:

(Don't forget to update vertex positions!)

# Continuity of loop subdivision surface 

- At extraordinary vertices
- Surface is at least ${ }^{1}$ continuous

■ Everywhere else ("ordinary" regions)

- Surface is ${ }^{2}$ continuous


## Loop subdivision results



# Catmull-Clark Subdivision 

## Catmull-Clark subdivision (regular quad mesh)



## Catmull-Clark subdivision (regular quad mesh)



## Catmull-Clark subdivision (regular quad mesh)



## Catmull-Clark vertex update rules (quad mesh)

Face point $\quad f=\frac{v_{1}+v_{2}+v_{3}+v_{4}}{4}$

$$
e=\frac{v_{1}+v_{2}+f_{1}+f_{2}}{4}
$$

Edge point



## Catmull-Clark subdivision (general mesh)



## Catmull-Clark subdivision (general mesh)



## Catmull-Clark subdivision (general mesh)



## Catmull-Clark subdivision (general mesh)



## Catmull-Clark vertex update rules (general mesh)

$f=$ average of surrounding vertices
$e=\frac{f_{1}+f_{2}+v_{1}+v_{2}}{4}$
These rules reduce to earlier quad rules for ordinary vertices/faces
$v=\frac{\bar{f}}{n}+\frac{2 \bar{m}}{n}+\frac{p(n-3)}{n}$
$\bar{m}=$ average of adjacent midpoints
$\bar{f}=$ average of adjacent face points
$n=$ valence of vertex
$p=$ old "vertex" point

## Continuity of Catmull-Clark surface

- At extraordinary points
- Surface is at least ${ }^{1}$ continuous
- Everywhere else ("ordinary" regions)
- Surface is ${ }^{2}$ continuous


## What about sharp creases?



## What about sharp creases?

Loop with Sharp Creases


Catmull-Clark with Sharp Creases


Figure from: Hakenberg et al. Volume Enclosed by Subdivision Surfaces with Sharp Creases

## Creases and boundaries

■ Can create creases in subdivision surfaces by marking certain edges as "sharp". Surface boundary edges can be handled the same way

- Use different subdivision rules for vertices along these "sharp" edges


Insert new midpoint vertex, weights as shown


Update existing vertices, weights as shown

## Subdivision in action ("Geri's Game", Pixar)

■ Subdivision used for entire character:

- Hands and head
- Clothing, tie, shoes



## Subdivision in action (Pixar's "Geri's Game")



## Mesh simplification (downsampling)

## How do we resample meshes? (reminder)

- Edge split is (local) upsampling:
- Edge collapse is (local) downsampling:

- Edge flip is (local) resampling:

- Still need to intelligently decide which edges to modify!


# Mesh simplification 

Goal: reduce number of mesh elements while maintaining overall shape


## Estimate: error introduced by collapsing an edge?

How much geometric error is introduced by collapsing an edge?


# Sketch of Quadric Error Mesh Simplification 

## Simplification via quadric error

- Iteratively collapse edges
- Which edges? Assign score with quadric error metric*
- Approximate distance to surface as sum of squared distances to planes containing nearby triangles
- Iteratively collapse edge with smallest score
- Greedy algorithm... great results!
* (Garland \& Heckbert 1997)


## Distance from point to a line (and a plane)

Line is defined by:

- Its normal: N
- A point $x_{0}$ on the line

$$
\begin{aligned}
& \mathbf{N} \cdot\left(\mathbf{x}-\mathbf{x}_{\mathbf{0}}\right)=0 \\
& \mathbf{N}^{\mathrm{T}}\left(\mathbf{x}-\mathbf{x}_{\mathbf{0}}\right)=0
\end{aligned}
$$

$$
\mathbf{N}^{\mathrm{T}} \mathbf{x}=\mathbf{N}^{\mathbf{T}} \mathbf{x}_{\mathbf{0}}
$$

$$
\mathbf{N}^{\mathrm{T}} \mathbf{x}=c
$$

Distance to line:
$\mathbf{N} \cdot\left(\mathbf{x}-\mathbf{x}_{\mathbf{0}}\right)$
The line (in 2D) is all points $x$, where x - $\mathrm{x}_{0}$ is orthogonal to N .
( $\mathrm{N}, \mathrm{x}, \mathrm{x}_{0}$ are 2 -vectors)

## Quadric error matrix (encodes squared distance)

- Suppose we have:
- a query point ( $x, y, z$ )
- a normal ( $\mathbf{a}, \mathrm{b}, \mathbf{c}$ )
- an offset $d:=-\left(x_{p}, y_{p}, z_{p}\right) \cdot(a, b, c)$

$$
Q=\left[\begin{array}{cccc}
a^{2} & a b & a c & a d \\
a b & b^{2} & b c & b d \\
a c & b c & c^{2} & c d \\
a d & b d & c d & d^{2}
\end{array}\right]
$$

- Then in homogeneous coordinates, let
- $\mathbf{u}:=(\mathbf{x}, \mathbf{y}, \mathbf{z}, \mathbf{1})$
- $v:=(a, b, c, d)$
- Signed distance to plane is then

$$
\mathbf{D}=\mathbf{u} \mathbf{v}^{\mathrm{T}}=\mathbf{v} \mathbf{u}^{\mathrm{T}}=\mathrm{ax}+\mathrm{by}+\mathrm{cz}+\mathbf{d}
$$

- Squared distance is $D^{2}=\left(u v^{\top}\right)\left(v u^{\top}\right)=u\left(v^{\top} v\right) u^{\top}:=u^{\top} \mathbf{Q u}$

- Distance is 2nd degree ("quadric") polynomial in $x, y, z$


## Cost of edge collapse

- How much does it cost to collapse an edge?
- Idea: compute edge midpoint $V_{\text {mid }}$, measure quadric error at this point
- Error at $V_{\text {mid }}$ given by $\mathbf{V}_{\text {mid }}{ }^{\top}\left(Q_{0}+Q_{1}\right) V_{\text {mid }}$ - See next slide for $Q_{i}$
- Intuition: cost is sum of squared differences to original position of triangles now touching $\mathbf{V}_{\text {mid }}$

- Better idea: choose point on edge (not necessarily the midpoint) that minimizes quadric error
- More details: Garland \& Heckbert 1997


## "Quadric error metric at mesh vertex"

Heuristic: "error metric at vertex $V$ " is sum of squared distances to triangles connected to $V$ Encode this as a single quadric matrix per vertex that is the sum of quadric error matrices for all triangles

$$
\mathrm{Q}_{\mathrm{V}}
$$



$$
Q_{V}=\sum_{i=1}^{N} Q_{i}
$$

## Quadric error simplification: algorithm

- Compute quadric error matrix Q for each triangle's plane
- Set Q at each vertex to sum of Q's from neighbor triangles
- Set Q at each edge to sum of Q's at endpoints
- Find point at each edge minimizing quadric error
- Until we reach target \# of triangles:

- collapse edge (i,j) with smallest cost to get new vertex m
- add $Q_{i}$ and $Q_{j}$ to get quadric $Q_{m}$ at vertex $m$
- update cost of edges touching vertex m



## Quadric error mesh simplification



## Mesh Regularization

## What makes a "good" triangle mesh?

■ One rule of thumb: triangle shape

- One rule of thumb: triangle shape
- More specific condition: Delaunay
- "Circumcircle interiors contain no vertices."
- Not always a good condition, but often*
- Good for simulation
- Not always best for shape approximation

*See Shewchuk, "What is a Good Linear Element"


## What else constitutes a good mesh?

■ Rule of thumb: regular vertex degree

- Triangle meshes: ideal is every vertex with valence 6:


Why? Better triangle shape, important for (e.g.) subdivision:

*See Shewchuk, "What is a Good Linear Element"

## Isotropic remeshing

Goal: try to make triangles uniform in shape and size


## How do we make a mesh "more delaunay"?

- Already have a good tool: edge flips!
- If $\alpha+\beta>\pi$, flip it!


■ In practice: a simple, effective way to improve mesh quality

## How do we improve degree?

- Edge flips!
- If total deviation from degree 6 gets smaller, flip it!


Iterative edge flipping acts like "discrete diffusion" of degree
No (known) guarantees; works well in practice

## How do we make triangles "more round"?

- Delaunay doesn't mean equilateral triangles
- Can often improve shape by centering vertices:



## Isotropic remeshing algorithm*

- Repeat four steps:
- Split edges over $4 / 3$ rds mean edge length
- Collapse edges less than $4 / 5$ ths mean edge length
- Flip edges to improve vertex degree
- Center vertices tangentially


## Things to remember

- Triangle mesh representations
- Triangles vs points+triangles
- Half-edge structure for mesh traversal and editing
- Geometry processing basics
- Local operations: flip, split, and collapse edges
- Upsampling by subdivision (Loop, Catmull-Clark)
- Downsampling by simplification (Quadric error)
- Regularization by isotropic remeshing


## Acknowledgements

- Thanks to Keenan Crane, Ren Ng, Pat Hanrahan, James 0'Brien, Steve Marschner for presentation resources

