Monte Carlo 2

Today

- Discrepancy and Quasi-Monte Carlo (QMC)
- Low-discrepancy constructions
 - Halton, Hammersley, Sobol'
- Randomized low-discrepancy: RQMC
- Spectral analysis of sampling patterns and MC

Warping Samples For MC Integration



$$\xi_i \in [0,1)^2$$

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- $\theta = 2\pi\xi_1$
- $r = \sqrt{\xi_2}$

Warping Samples For MC Integration

Cosine-weighted hemisphere sampling:



Three 2D Point Sets



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Point Set Evaluation: Discrepancy





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 $\Delta(x, y) = \frac{n(x, y)}{N} - xy$

 $D_N = \max_{x,y} \left| \Delta(x,y) \right|$

Three 2D Point Sets



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Discrepancy (Empirical)



IndependentStratified0.1480.081

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Larscher-Pillichshammer 0.041

Low-Discrepancy Definition

An (infinite) sequence of n samples in dimension d is low discrepancy if:

$$D_n = O\left(\frac{(\log n)^d}{n}\right)$$

A (finite) set of n samples in dimension d is low discrepancy if:

$$D_n = O\left(\frac{(\log n)^{d-1}}{n}\right)$$

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Theorem on Total Variation

Koksma-Hlawka inequality:

 $\left|\frac{1}{N}\sum_{i=0}^{N-1}f(X_i) - \int f(x)dx\right| \leq V(f)D_N$

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 $V(f) = \int \left| \frac{\delta f}{\delta x} \right| \, dx$

Quasi-Monte Carlo Error Bounds

Although error is bounded as $|e| \leq V(f)D_N$ not a tight bound!

Even worse, V(f) is sometimes unbounded

We can use this inequality to show that QMC error converges as:

$$\sim \frac{(\log N)^d}{N}$$

(recall that MC variance goes at O(1/N), so error goes at $O(1/\sqrt{N})$.)

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Measured Quasi-Monte Carlo Error



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- Independent
- Stratified
- Halton

- Independent
- Stratified
- Halton

Low-Discrepancy Point Sequences

The Radical Inverse

Consider the digits of a number n, expressed in base b $n = \sum d \cdot h^{(i-1)}$

$$i = \sum_{i=1}^{n} d_i b^{(i-1)}$$

e.g. for n = 6 in base 2, $n=110_2$, and $d_1 = 0, d_2 = 1, d_3 = 1, d_i = 0$

The radical inverse mirrors the digits around the decimal:

 $\Phi_2(6) = 0.011_2 = 0 \cdot 2^{-1} + 1 \cdot 2^{-2} + 1 \cdot 2^{-3} = 0.375$

$$\Phi_b(n) = \sum_{i=1}^{\infty} d_i \, b^{-i}$$

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1D Low Discrepancy: van der Corput

n	$\Phi_2(n)$
0	0
1	0.5
2	0.25
3	0.75
4	0.125
5	0.625
6	0.375
7	0.875
•••	•••



Efficient Base 2 Radical Inverse

Assume a fixed number of bits (say 32):

$$\Phi_b(n) = \sum_{i=1}^{52} d_i \, b^{-i}$$

29

We have the sum: $d_1 2^{-1} + d_2 2^{-2} + \cdots + d_{32} 2^{-32}$

Pull out a factor of 2^{-32} : $2^{-32}(d_1 2^{31} + d_2 2^{30} + \cdots + d_{32})$

Can also express in terms of bit shifts:

 $2^{-32}((d_1 << 31) + (d_2 << 30) + \dots + d_{32})$

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Efficient Base 2 Radical Inverse

$$2^{-32}((d_1 << 31) + (d_2 << 30)$$

We already have the digits in the bits of n

$$n = \sum_{i=1}^{\infty} d_i b^{(i-1)}$$

32	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1

So

Reverse the bits Multiply by 2^{-32}

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$+\cdots + d_{32})$

Reversing Bits in Parallel

ui

nt32	2 t	: Re	vei	rseI	Bit	s (u	int	32 1	tn)	{		32	31 30	29	28	27	26	25	24	23	22	21	20	19	18
		1		1.0				1.0		•															
n	=	(n	<<	10)		(п	. >>	το));			16	15 14	13	12	11	10	9	8	7	6	5	4	3	2
												16	15 14	13	12	11	10	9	8	7	6	5	4	3	2
n	=	((n	. &	0x0)0f	£00	ff)	<<	8)	I	((n	& 8	0xf 7 6	f0	0f 4	£0 3	0) 2	> 1	>> 16	8 15); 14	13	12	11	1(
												8	7 6	5	4	3	2	1	16	15	14	13	12	11	1(
																	•								
n	=	((n	L &	0x0)±0	±0±	: 0 £)	<<	4)	I	((n	& 4				0 ±	0) 6	>	•>	4);	9	16	15	14
												4	3 2	1	8	7	6	5	12	11	10	9	16	15	14
n	=	((n	. &	0x 3	333	333	33)	<<	2)	I	((n	&	0xc	CC	cc	cc	c)	>	·>	2));				
												2	1 4	3	6	5	8	7	10	9	12	11	14	13	16
												2	1 4	3	6	5	8	7	10	9	12	11	14	13	16
n re	= etu	((n 	n;	0x5	555	555	55)	<<	1)	Ι	((n	&	0xa	aa	aa	aa	a)	>	>	1);				
												1	2 3	4	5	6	7	8	9	10	11	12	13	14	15



Efficient van Der Corput

```
uint32 t ReverseBits(uint32 t n) {
  n = (n << 16) | (n >> 16);
   n = ((n \& 0x00ff00ff) << 8) | ((n \& 0xff00ff00) >> 8);
  n = ((n \& 0x0f0f0f0f) << 4) | ((n \& 0xf0f0f0f0) >> 4);
   n = ((n \& 0x3333333) << 2) | ((n \& 0xccccccc) >> 2);
   n = ((n \& 0x5555555) << 1) | ((n \& 0xaaaaaaaa) >> 1);
   return n;
```

```
float RadicalInverse2(uint32 t v) {
v = ReverseBits(v);
const float Inv2To32 = 1.f / (1ull << 32);
return v * Inv2To32;
```

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Radical Inverse Base 3

n	$\Phi_3(n)$
0	0
	0.333
2	0.666
3	0.111
4	0.444
5	0.777
6	0.222
7	0.555
8	0.888



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The Halton Sequence

Low discrepancy sequence $(\Phi_{b_1}(n), \Phi_{b_2}(n), \Phi_{b_3}(n), \ldots)$ The dimensions' bases are relatively prime Arbitrary number of dimensions Arbitrary number of points



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The Hammersley Point Set

If the number of points N is known in advance, set one dimension to n/N



Slightly lower discrepancy than Halton

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Low-Dimensional Projections...

Caution: 2D projections of higher bases may not be great

The overall pattern remains low-discrepancy over all dimensions, though



 $(\Phi_{29}(n), \Phi_{31}(n))$

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Randomized Low Discrepancy

$$\Phi_b(n) = \sum_{i=1}^{\infty} d_i b^{-i} \qquad \Phi'_b(n)$$

Radical Inverse

σ_i are (random) permutations of the digits Random permutations maintain LD

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 $=\sum \sigma_i(d_i)b^{-i}$ i=1

Permuted Radical Inverse

Halton + Random Digit Permutations



Unscrambled

 $(\Phi_{29}(n), \Phi_{31}(n))$

Scrambled

Owen Scrambling

Apply random digit permutations that depend on previous digits

$$\Phi'_b(n) = \sum_{i=1}^{\infty} \sigma_i(d_i) b^{-i}$$

Permuted Radical Inverse

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Error With Owen Scrambling



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- Independent
- Sobol' (Owen)
- Sobol' (Digit Permute)

- Independent
- Sobol' (Owen)
- Sobol' (Digit Permute)

Sobol' Point Sets



Generator Matrices

Given a base b and a matrix C, define:

• where d_i are the base-b digits of n \blacksquare and arithmetic is done over the ring \mathbb{Z}_b For our purposes, just do everything "mod b" This generates a set of b^m points



Generator Matrices

We'll focus only on b=2, which allows particularly efficient implementation

$$c(n) = (2^{-1}, 2^{-2}, \dots, 2^{-m})$$









Sobol' Point Sets

Sobol' showed how to find generator matrices for LD point sets in base 2

Can scale low-discrepancy samples in 1000s of dimensions



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 $\bullet \bullet \bullet$



32 2D Sobol' Points



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Elementary Intervals (1x32)



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Elementary Intervals (2x16)



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Elementary Intervals (4x8)



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Elementary Intervals (8x4)



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Elementary Intervals (16x2)



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Elementary Intervals (32x1)



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Independent Random Samples, n=16 MSE 1x



Stratified Samples, n=16 MSE 1/2.41x



Sobol' Samples, n=16 MSE 1/3.38x



Independent Random Samples, n=16 MSE 1x



Stratified Samples, n=16 MSE 1/2.12x



Sobol' Samples, n=16 MSE 1/3.95x

Warping Samples to a Quad Light



4x4 Stratified



16 Sobol'

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Sampling Motion Blur + Defocus







Halton MSE 1/1.13x

Independent MSE 1x

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Sobol' MSE 1/1.80x

(0,2)-sequences

In addition to satisfying general stratification properties, power-of-two length subsequences are well-distributed with respect to each other





Pixel * Light Sampling



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Spectral Analysis of Sampling

Measuring Point Set Quality

- Some problems with discrepancy:
 - Anisotropic: rotating the points changes discrepancy
 - Not shift-invariant: similarly for translation
 - Doesn't account for human perception
- In general, can have low discrepancy yet still have points clumped together:

Ambient Occlusion: $\int_{\Omega} V(\omega) \cos \theta \, d\omega$





Reference

Random A

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Random B









Blue Noise Dithering (Ulichney)





White

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Blue

Power Spectrum of Samples

Samples



Fourier Transform



 $P_f(\omega) = F(\omega)F(\omega)$ $= F(\omega)^2$

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Radial Power Spectrum



is even

Colors of Noise

White



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ER SPECTRUM

РОМ

Ambient Occlusion, Revisited $\int_{\Omega} V(\omega) \cos \theta \, d\omega$





Reference

Blue Noise

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White Noise

(0,2)-sequence









Blue Noise Sobol'





Jittered Power Spectrum



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Low Discrepancy Power Spectra



Halton

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Power Spectra and Aliasing







Halton



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Randomized LD Power Spectra



Sobol'

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Owen Scrambled Sobol'

Integration Error when Sampling

Integral

$$I(f) = \int f(x) \, dx$$

Sampled integral

Given samples x_i , define $s(x) = \frac{1}{N} \sum_{i=1}^{N} \delta(x - x_i)$

$$I_s(f) = \int f(x)s(x) \, dx = \frac{1}{N} \sum f(x_i)$$

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Integration Error when Sampling

Sampled integral

$$I_s(f) = \int f(x)s(x) \, dx = \frac{1}{N} \sum f(x_i)$$

Since:

$$f(x)s(x) \leftrightarrow F(\omega) \oplus S(\omega)$$

$$I(f) = \int f(x) \, dx = F(0)$$

$$I_s(f) = \int f(x)s(x) \, dx = F(\omega) \oplus S(\omega)$$

Error

 $\Delta = F(0) - F(\omega) \oplus S(\omega)|_{\omega=0}$

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Integration in the Frequency Domain



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[Durand 2011]

Variance Analysis

Recall error is: $\Delta = F(0) - F(\omega) \oplus S(\omega)|_{\omega=0}$

Can show that: $V \propto \int_{\Omega - \{0\}} P_F(\omega) P_S(\omega) \,\mathrm{d}\omega$

where $P_*(\omega)$ is the power spectrum of the signal

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Power Spectra and Variance



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Lecture 18

[Pilleboue et al. 2015]